Course Introduction

and some motivation

Mainack Mondal

CS60112
Spring 2025

Today’s Class

- Course logistics
- How to learn security?
- How to (not) nuke your system?
- Epilogue

Instructor

Mainack Mondal: usable security and privacy, system security and privacy, operationalizing privacy theories

Office: CSE 316

TA’s

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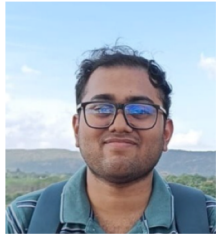

Sourabh Soumyakanta Das

Dassourabh103 at
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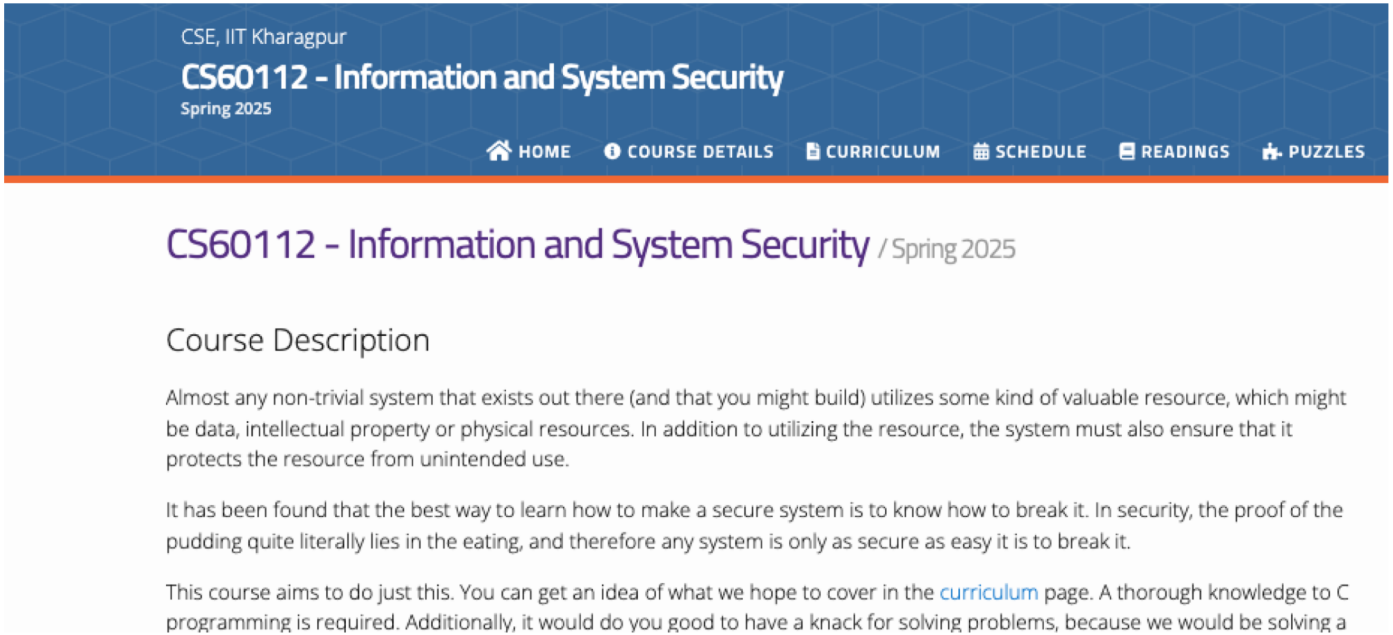
Website

https://kronos-192081.github.io/InfoSec-2025/



Course Timings

- Credit: 3-0-0

- Wednesday 10:00 am - 10:55 am
- Thursday 9:00 am - 9:55 am
- Friday 11:00 am - 11:55 am

Mode of Teaching

- Offline lectures
  - Please come to class (no recordings)
- (occasional) Pre-recorded lectures for special topics
  - We will upload the recorded lectures via MS Teams
- Two exams + Assignments

MS Teams

Link:

https://teams.microsoft.com/l/team/19%3AlF59SxbMNlF8qcEtIaabtg4wZyRFj8kfPVO_yfyZftU1%40thread.tacv2/conversations?groupId=86e1aa23-cbd8-477b-9b38-2d6f7b6f5dec&tenantId=71dbb522-5704-4537-9f25-6ad2dcd4278d

Or

https://shorturl.at/5UUW7

Code: **m6axutv**

Course evaluation: Exam

Two exams (55%)

- Syllabus: Everything until that point
- Dates will be in the webpage and announced in academic calendar

Course evaluation: Assignments

Regular assignments (45%)

- You will learn this course mostly through the assignments
- Assignments will comprise CTF style problems
- Start forming teams! (size <= 3)

Course logistics

Questions?

Ethical Considerations

Source: https://mozzoneplayer.com/2014/11/22/the-4th-monkey-do-no-evil/

Ethical Considerations

- Don’t do evil
- If you feel it is wrong, it is wrong
- Cyber offenses are punishable by law
- Use your tools responsibly (“It was an accident, Milord!” won’t hold up)

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How to learn Security?

You first learn how to break systems!

How to learn Security?

Then you’d know what to not do!

CTF

Capture the Flag (CTF) is a kind of security exercise, where the aim is to “break” the system in order to retrieve a piece of text, called the “flag”.

CTF rules

- The flag must be submitted **as-is** (no modifications)
- Checking will be **automated**
- Flags will be **randomized** on a per-team basis
- Sometimes the flag will have a specified format, such as “drapeau{<flag-text>}”. You need to submit the **whole thing**.

Leaderboard

- We will have a leaderboard for assignments.
- All points add up.
- Top three teams to get cash prizes!

Topics

- Web security
- Reverse Engineering
- Pwning
- Cryptography
- Hardware

Topic 1: Web security

- Perhaps the source tells you something
- Perhaps you can create a url which will lead you to a secret
- …
- Goal: find vulnerabilities

Topic 2: Reverse engineering

- Auditing binaries (static, dynamic)
- Understand the binary code / file
- ...
- Goal: check if you can find vulnerabilities using that knowledge

Topic 3: Pwning

- Pwning -> **O**wning
- Find and exploit vulnerabilities to obtain access to a system

Topic 3: Pwning

- Demo?

Topic 3: Pwning

- https://www.programiz.com/c-programming/online-compiler/

```
#include <stdio.h>
int main() {
    char *buf = "iss2024";
    puts(buf);
    return 0;
}
```

What would it print?

Topic 3: Pwning

- https://www.programiz.com/c-programming/online-compiler/

```
#include <stdio.h>
int main() {
    char *buf = "iss2024";
    puts(buf);
    return 0;
}
```

What would it print? **Make it a bit more interesting?**

Topic 4 / 5: Cryptography / HW

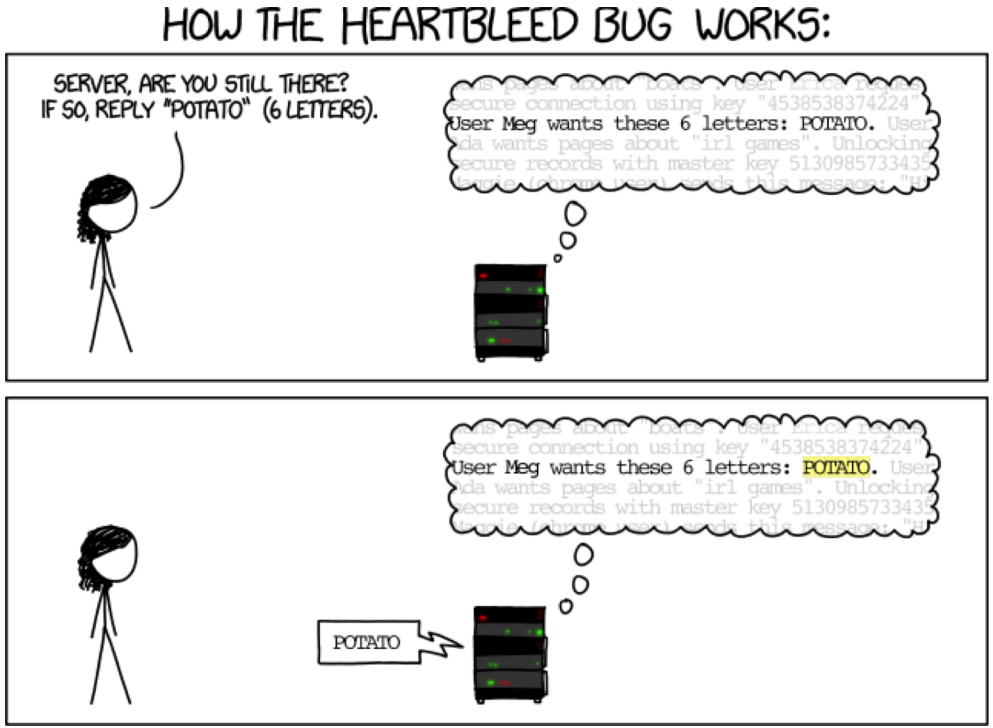
- Will cross the bridge when we get there

Topic 4 / 5: Cryptography / HW

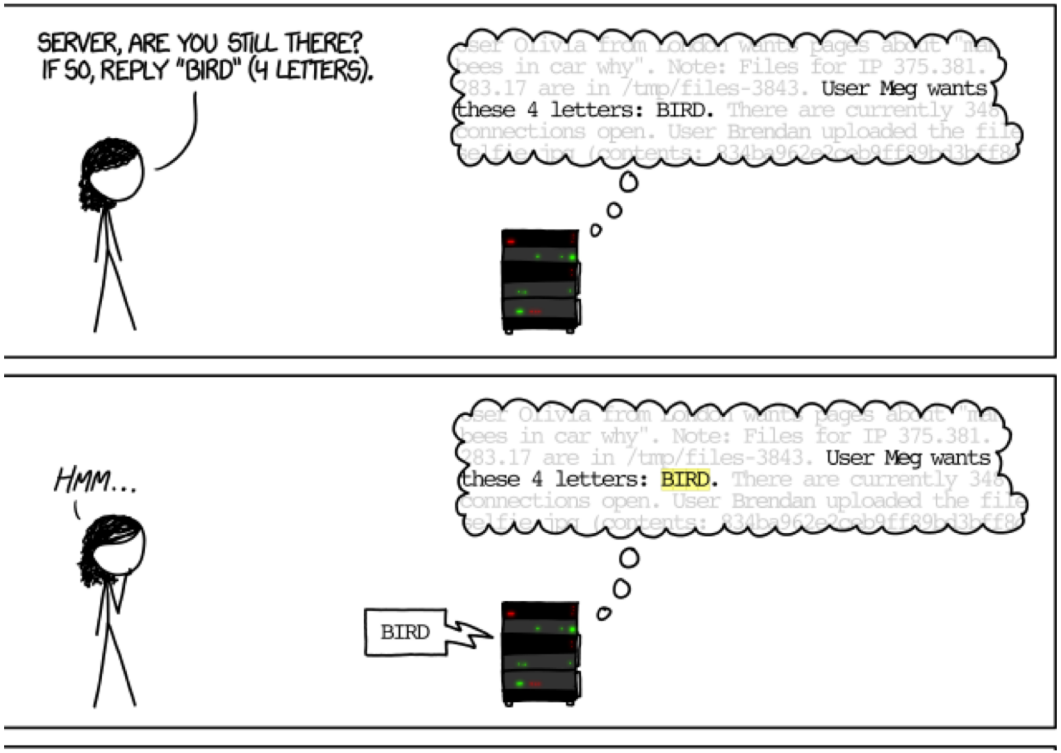
break

- Will ~~cross~~ the bridge when we get there

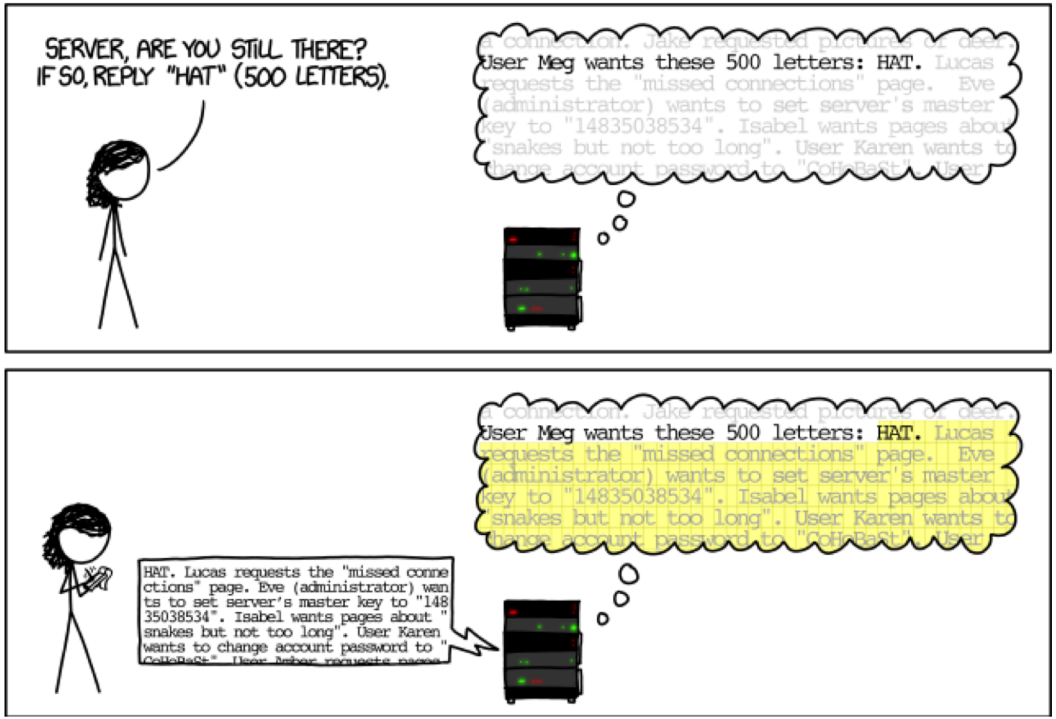
Case study 1: Heartbleed (1/3)



Case study 1: Heartbleed (2/3)



Case study 1: Heartbleed (3/3)



```
1448 dtls1_process_heartbeat(SSL *s)
1449   {
1450   unsigned char *p = &s->s3->rrec.data[0], *pl;
1451   unsigned short hbtype;
1452   unsigned int payload;
1453   unsigned int padding = 16; /* Use minimum padding */
1454
1455   /* Read type and payload length first */
1456   hbtype = *p++;
1457   n2s(p, payload);
1458   pl = p;
...
1465   if (hbtype == TLS1_HB_REQUEST)
1466     {
1467     unsigned char *buffer, *bp;
...
1470     /* Allocate memory for the response, size is 1 byte
1471     * message type, plus 2 bytes payload, plus
1472     * payload, plus padding
1473     */
1474     buffer = OPENSSL_malloc(1 + 2 + payload + padding);
1475     bp = buffer;
1476
1477     /* Enter response type, length and copy payload */
1478     *bp++ = TLS1_HB_RESPONSE;
1479     s2n(payload, bp);
1480     memcpy(bp, pl, payload);
```

(a) The Heartbeat buggy C code in `ssl\d1_both.c` [10].

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(a) The Heartbeat buggy C code in `ssl\d1_both.c` [10].

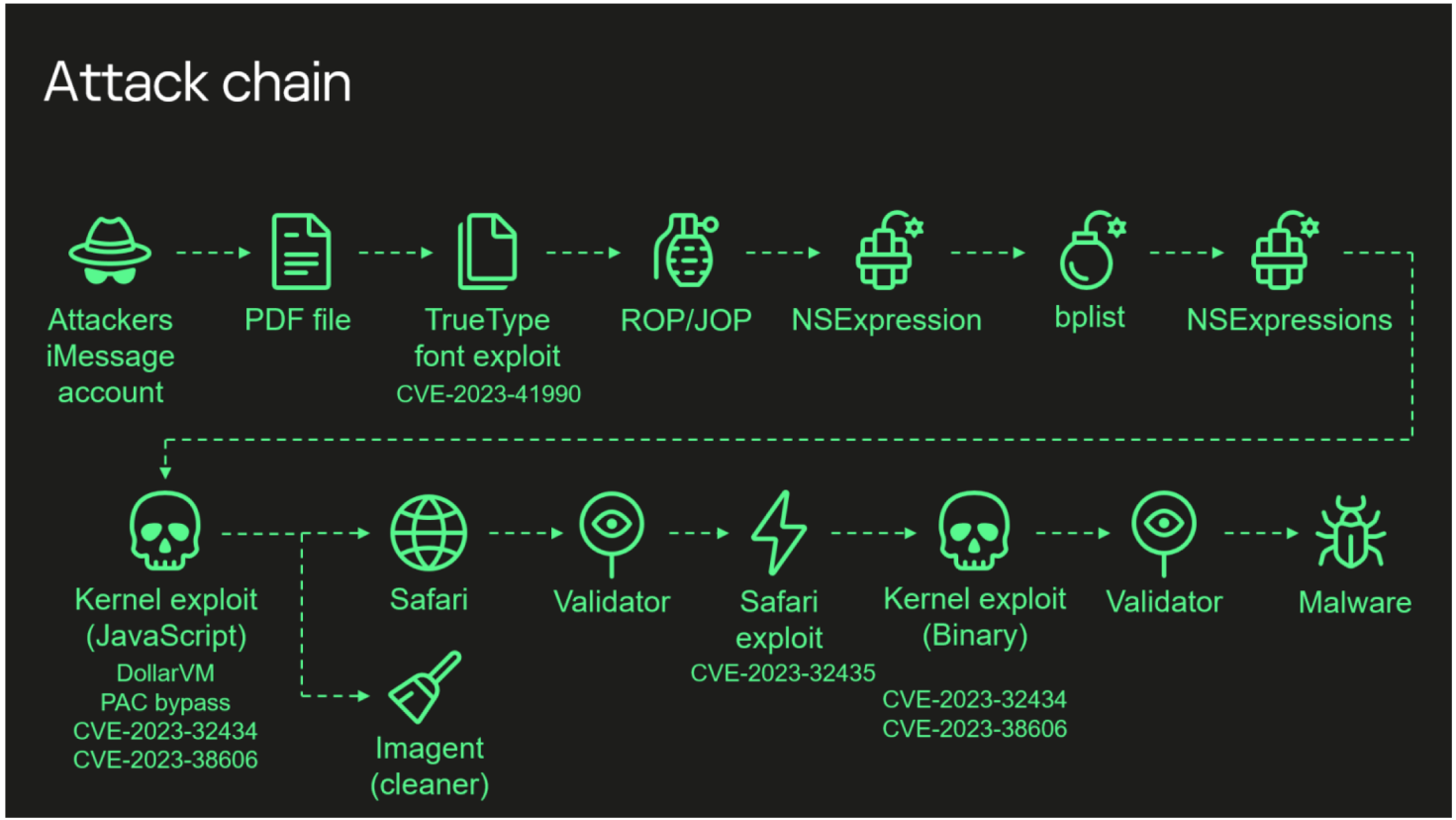
```
/* Naive implementation of memcpy
void *memcpy (void *dst, const void *src, size_t n)
{
    size_t i;
    for (i=0; i<n; i++)
        *(char *) dst++ = *(char *) src++;
    return dst;
}
```

payload, bp, pl

(b) A naive C implementation of the `memcpy()` function.

FIGURE 1: Analysis of Heartbleed.

Case study 2: Operation Triangulation (1/2)



Case study 2: Operation Triangulation (2/2)

However apple had hardware protection against writing random memory addresses

Case study 2: Operation Triangulation (2/2)

However apple had hardware protection

- Against writing random memory addresses

Bypassed using undocumented registers

- possibly there for debugging

https://securelist.com/operation-triangulation-the-last-hardware-mystery/111669/

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How to nuke your system?

- Run untrusted executables/programs on your system
- Turn off kernel security features
- Stick “sudo” before random commands
- Run “sudo rm -rf --no-preserve-root /”

How to NOT nuke your system?

Ignore all the previous advice and instead:

- Use Docker (has caveats)
- Use virtualization
  - VirtualBox
  - QEMU

Docker

- Lightweight “virtualization” software, with the caveat being it requires root to run (so be careful with mounted volumes).
- Kernel shared with host with namespace based isolation.
- Install Docker Engine (required).

VirtualBox

- Full-system virtualization software, with a nice GUI.
- Install from here (required).
- Also install the extension pack and guest additions for nicer integration.
- We will provide VirtualBox VM files (with extension .ovf) when required.

QEMU

- Full-system virtualization + architecture emulation software, CLI only.
- Install from here (required).
- QEMU typically requires a long list of flags and a disk-image file to run. We will provide both.
- Helpful in emulating non-native ISAs (e.g., MIPS, RISC-V)

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